



Equivalence of Uniform Key Agreement and Composition Insecurity

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Plan

- Main Question:
 - hardness amplification via composition
- Previous Works
- Our main result
- Conclusions



Types of Compositions

- **Parallel Composition**
Let F and G be functions with the same domain and range. Parallel composition of F and G is defined as $P(\cdot)$:

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- **Sequential Composition**

Let F and G be functions with the same domain and range. Sequential composition of F and G is defined as $S(\cdot)$:

$$S(\cdot) = G(F(\cdot))$$



Security Amplification via Composition

- Does **composition** of Non-adaptively secure PRF (unconditionally) implies adaptive security?
 - - In an information-theoretic setting?
 - - In a computational setting?
 - If not, under what assumptions?



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 - - In a computational setting?
 - If not, under what assumptions?
- A lot is known already, lets review...
[LR86], [Vaud03], [MP 04], [Piet05], [Piet06]



Previous Works

- Information Theoretic Setting
 - **Vaudenay [Vaud 03]** showed the **sequential** composition of k non-adaptively ϵ -secure Permutation implies $2^{k-1} \epsilon^k$ adaptive security.



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- **Information Theoretic Setting**
 - **Vaudenay [Vaud 03]** showed the **sequential** composition of k non-adaptively ϵ -secure Permutation implies $2^{k-1} \epsilon^k$ adaptive security.
 - **Luby and Rackoff [LR86]** showed the **parallel** composition of k ϵ -secure functions implies $2^{k-1} \epsilon^k$ adaptive security.



Previous Works (Cont)

- Information Theoretic Setting
- **Maurer and Pietrzak [MP 04]** showed for permutations that composition of 2 ε -secure non-adaptive permutations implies $2\varepsilon(1 + \ln(\varepsilon^{-1}))$ secure adaptive permutation.



Previous Works (Cont)

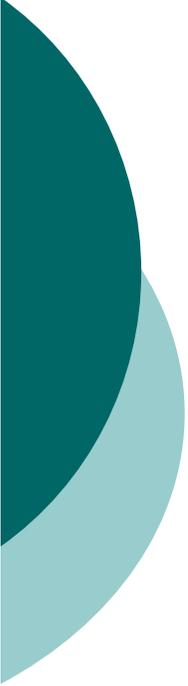
- Information Theoretic Setting
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- Computational Setting
 - **Pietrzak [Piet05]** showed:
 - **DDH** \rightarrow composition does not help, i.e.:
 - If **DDH** assumption holds, then the composition of non-adaptively secure functions is **not** adaptively secure (i.e. 3 adaptive-query breakable)!



Previous Works (Cont.)

- **Pietrzak [Piet06]:**

- *If sequential composition does not help (for PRF) → there exists Key-agreement*
- If the *sequential* composition of k -adaptively secure functions is not $k+1$ adaptively secure, then there exists a $(2k-1)$ -pass key agreement.
- In fact, the above construction of key agreement implies **uniform transcript key agreement**.
- *If sequential composition does not help → uniform-transcript Key-agreement*



Summary of [Piet05] and [Piet06]:

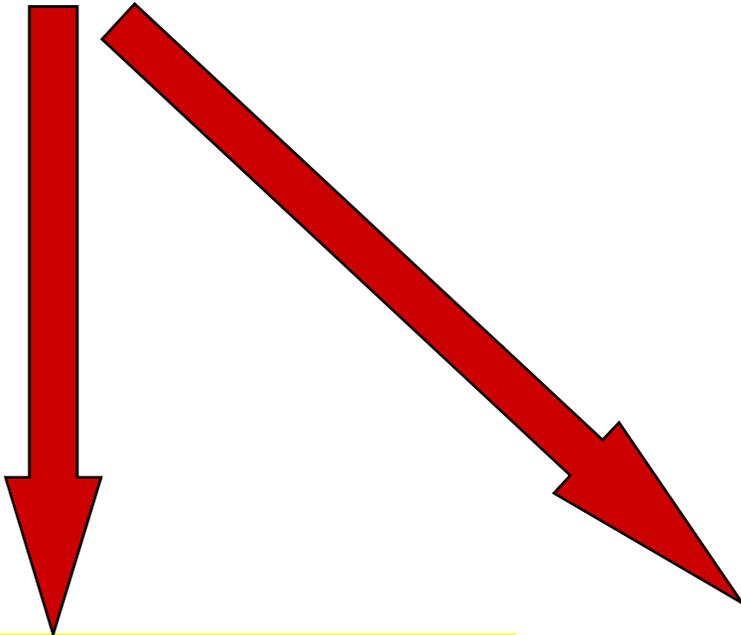
DDH

Parallel composition
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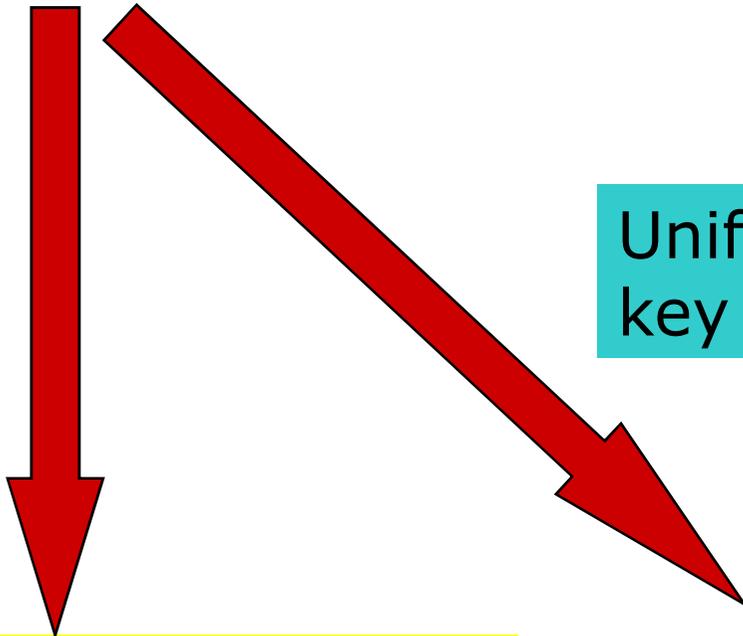
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Uniform transcript
key agreement

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What else do we know?

DDH

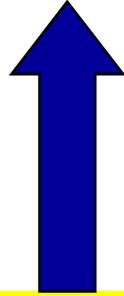
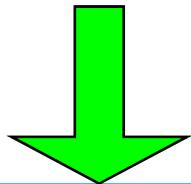
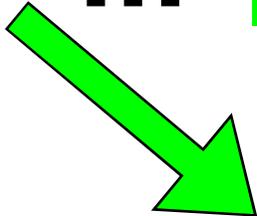
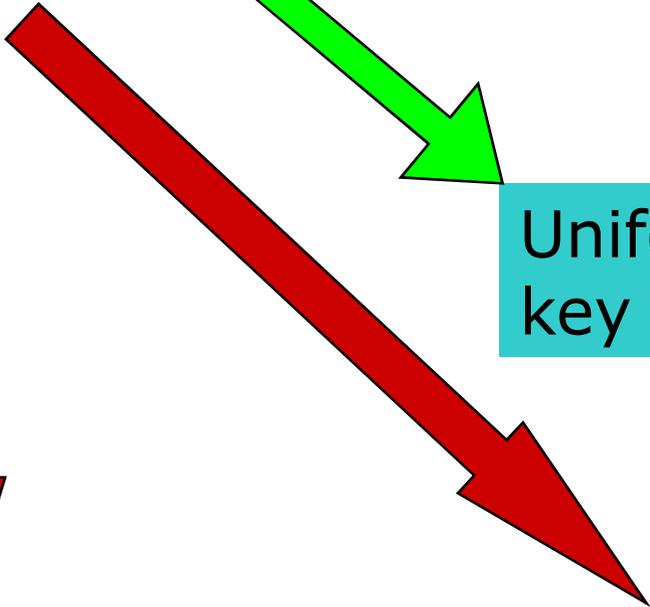
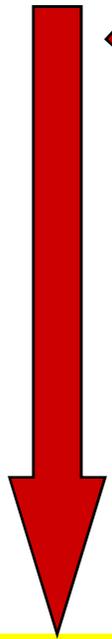
...

Dense trapdoor permutations

Uniform transcript
key agreement

Parallel composition
does not help

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Wishful thinking...

DDH

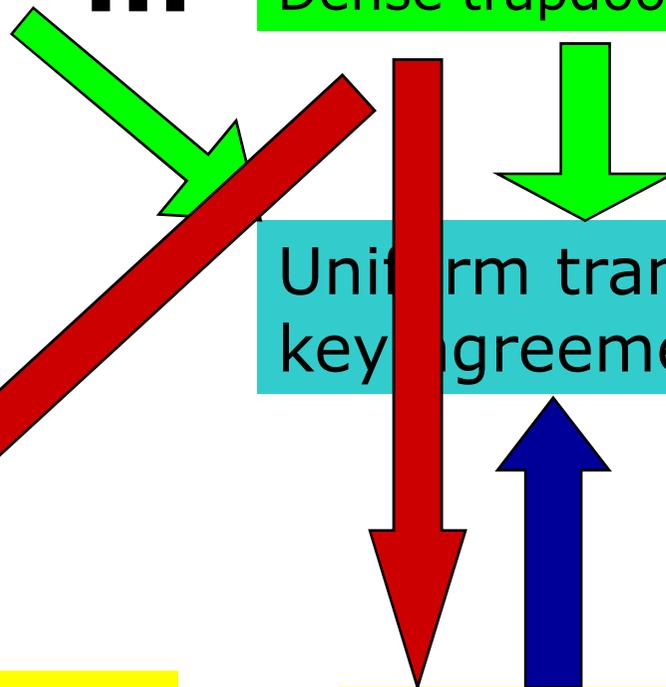
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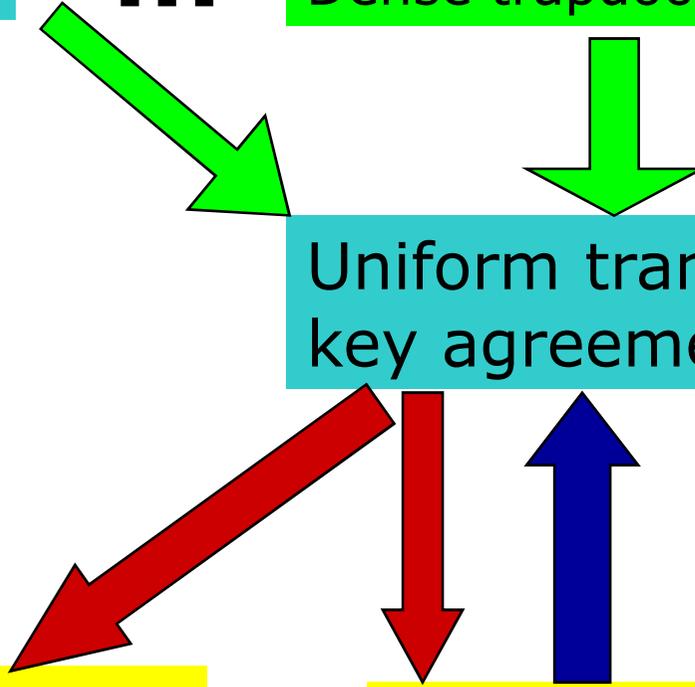
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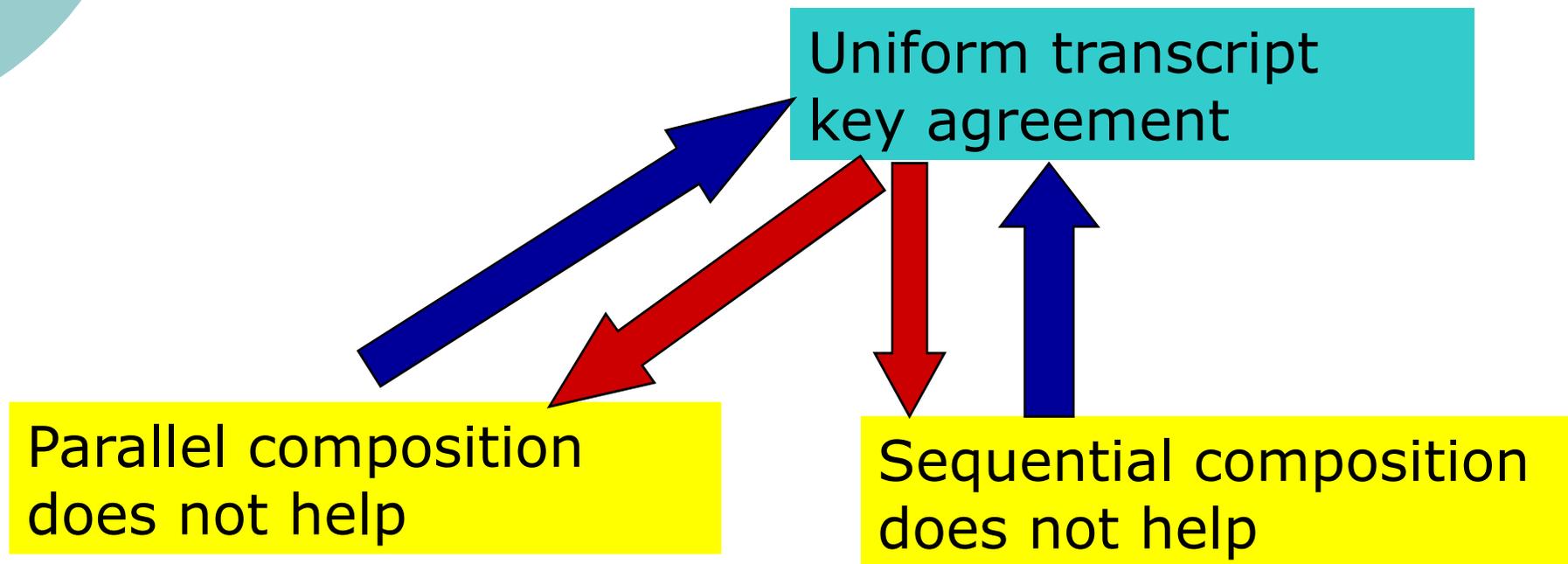
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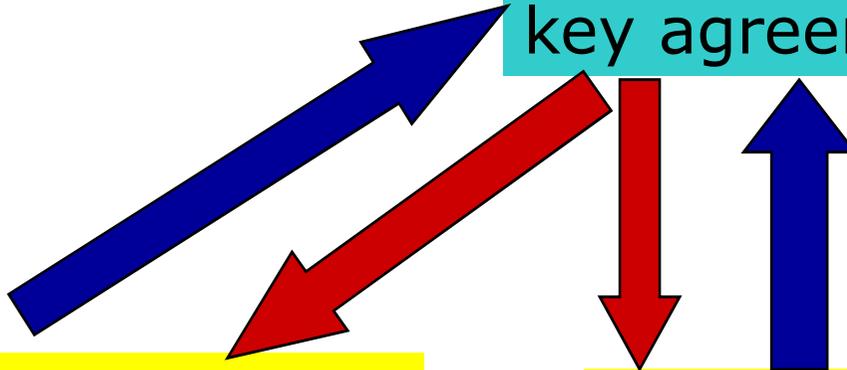
Public Key World



Uniform transcript
key agreement

Parallel composition
does not help

Sequential composition
does not help



THIS IS EXACTLY WHAT WE DO

Public Key World



Uniform transcript
key agreement

Parallel composition
does not help



Sequential composition
does not help

Private Key World

THIS IS EXACTLY WHAT WE DO

Public Key World



Uniform transcript
key agreement



Parallel comp
does not help

composition
elp

Private Key World



Our Result

- **Main Thm: Both sequential and parallel composition of two pseudo-random functions does not imply adaptive-security**

if and only if

a uniform-transcript key agreement exists.



Conclusion and Open Questions

- Composition Insecurity \leftrightarrow Uniform Transcript Key Agreement
- Round Complexity of UTKA is Linearly Proportional to the Adaptive Security of Compositions.
- Open: Tighter Relation between the security of component functions and the security of their compositions.