

# Protecting Circuits from Computationally-Bounded Leakage

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# Motivation

**The great tragedy of Crypto –  
the slaying of a provably secure scheme  
by an ugly side channel.**

# Engineering approach

Try preventing side-channel leakage.



# Engineering approach

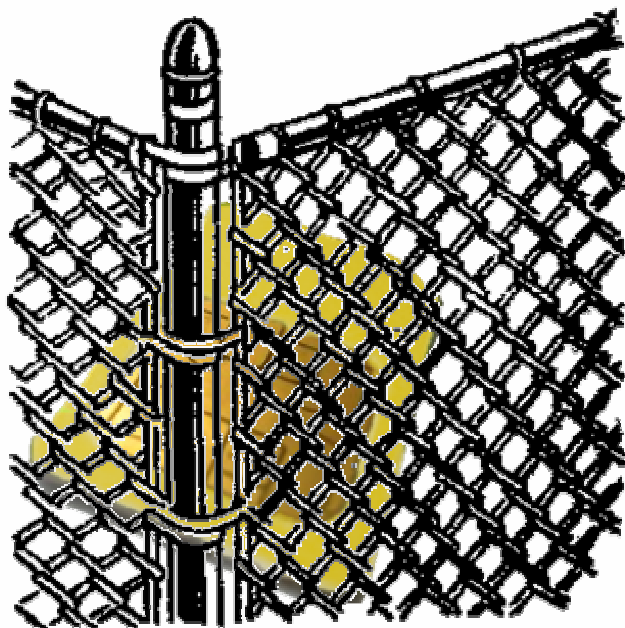
Try preventing side-channel leakage.





# Engineering approach

Try preventing side-channel leakage.



# Engineering approach

Try preventing side-channel leakage.



Good luck...



# Cryptographic approach

Leakage is a *given*,  
modeled by an adversarial  
observer.

The device should  
protect itself against it.



## Related Work

[CDHKS00]: Canetti, Dodis, Halevi, Kushilevitz, Sahai: Exposure-Resilient Functions and All-Or-Nothing Transforms

[ISW03]: Ishai, Sahai, Wagner: Private Circuits: Securing Hardware against Probing Attacks

[MR04]: Micali, Reyzin: Physically Observable Cryptography

[GTR08]: Goldwasser, Tauman-Kalai, Rothblum: One-Time Programs

[DP08]: Dziembowski, Pietrzak: Leakage-Resilient Cryptography in the Standard Model

[Pie09]: Pietrzak: A leakage-resilient mode of operation

[AGV09]: Akavia, Goldwasser, Vaikuntanathan: Simultaneous Hardcore Bits and Cryptography against Memory Attacks

[ADW09]: Alwen, Dodis, Wichs: Leakage-Resilient Public-Key Cryptography in the Bounded Retrieval Model

[FKPR09]: Faust, Kiltz, Pietrzak, Rothblum: Leakage-Resilient Signatures

[DHT09]: Dodis, Lovett, Tauman-Kalai: On Cryptography with Auxiliary Input

[SMY09]: Standaert, Malkin, Yung: A Unified Framework for the Analysis of Side-Channel Key-Recovery Attacks

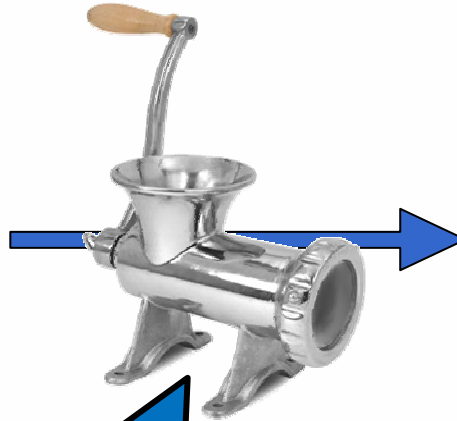
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# Model

[Ishai Sahai Wagner '03]



**Any** boolean circuit



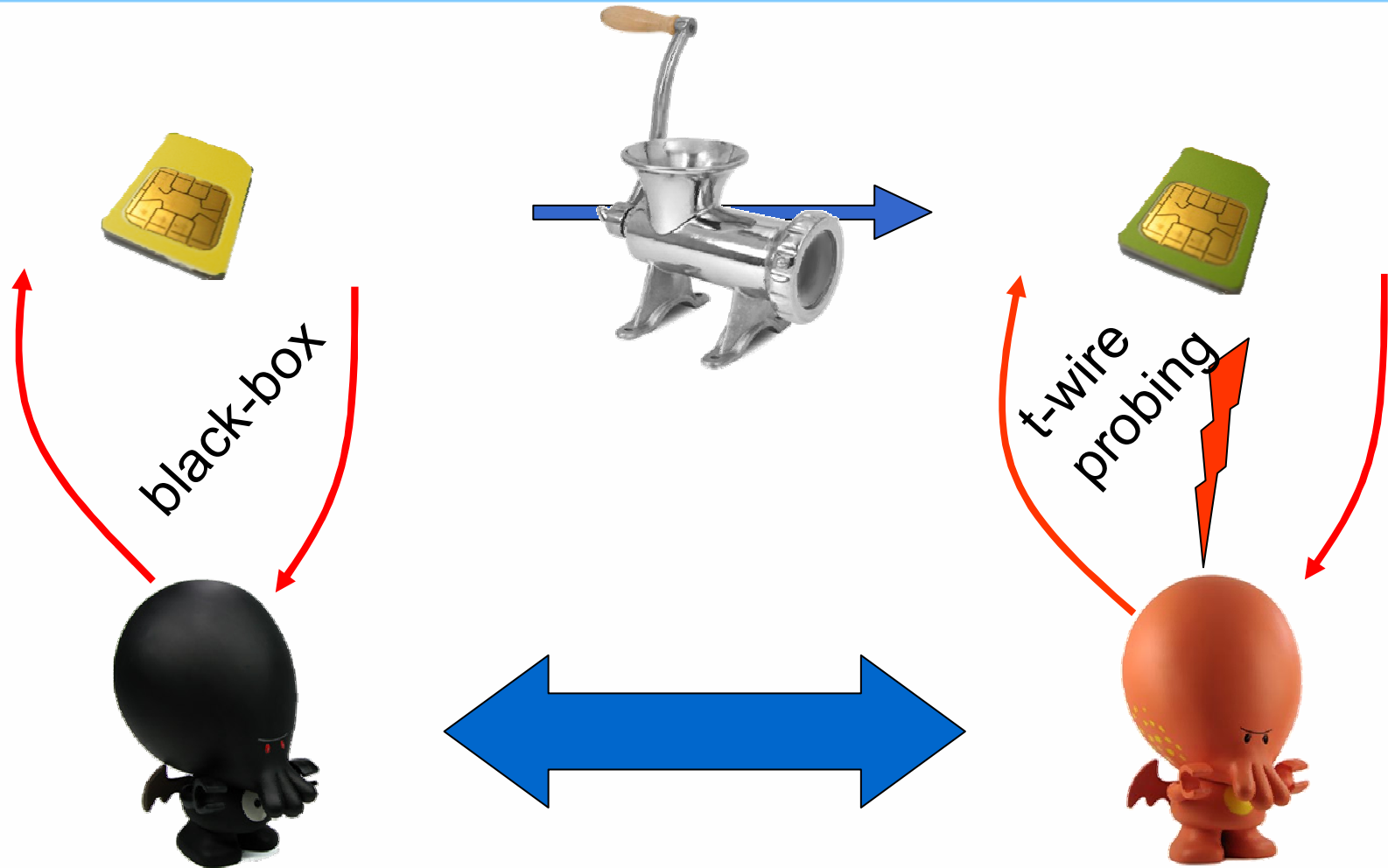
Circuit  
transformation



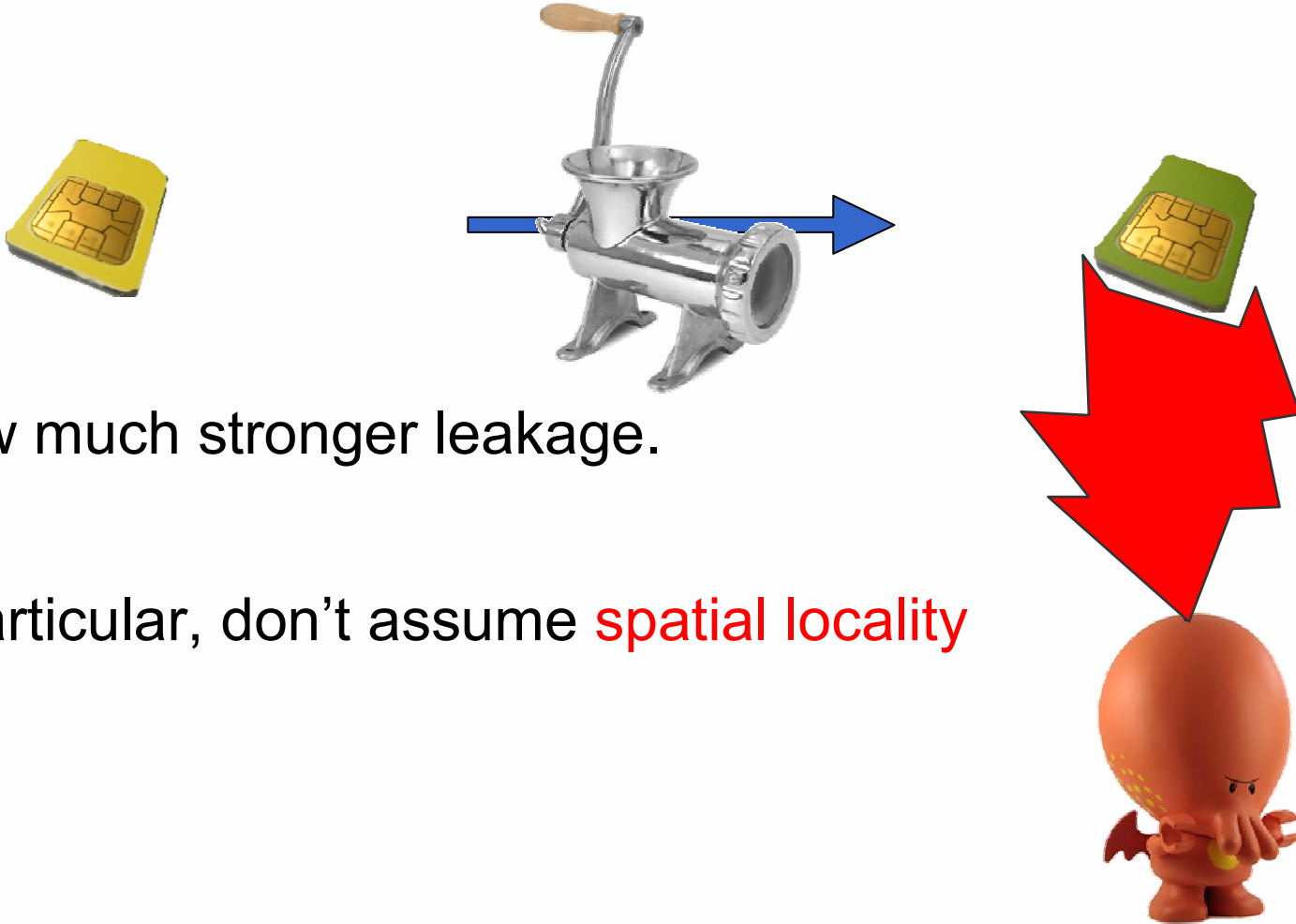
Transformed circuit

# Model

[Ishai Sahai Wagner '03]



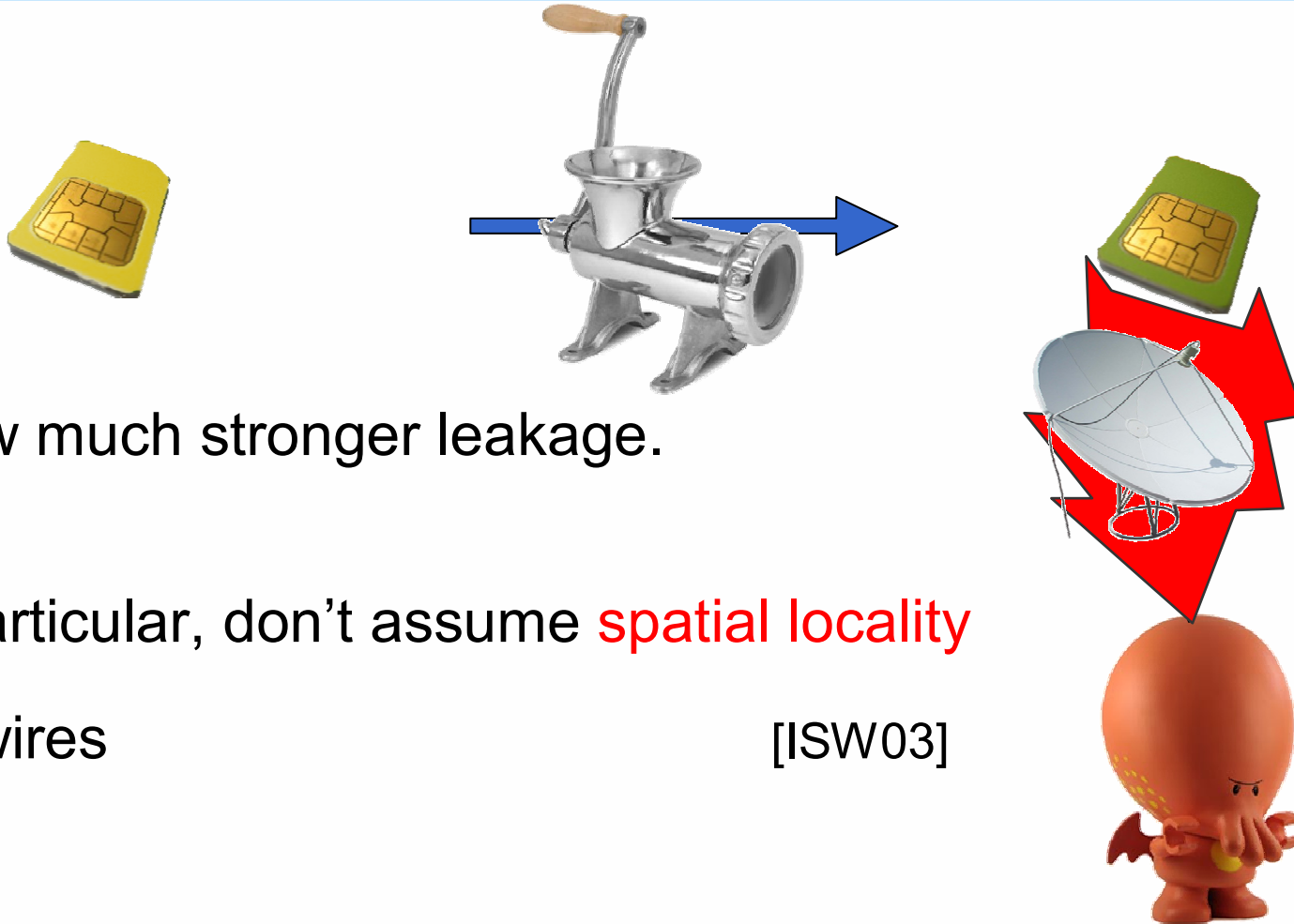
# Our goal



Allow much stronger leakage.

In particular, don't assume **spatial locality**

# Our goal



Allow much stronger leakage.

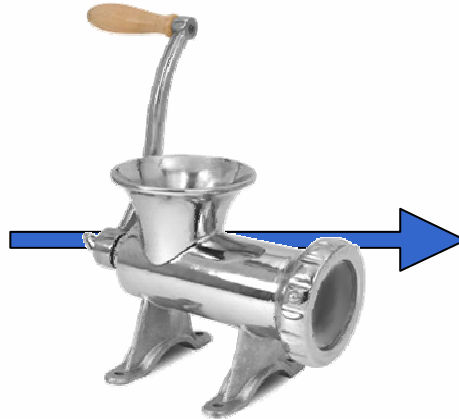
In particular, don't assume **spatial locality**

- $t$  wires

[ISW03]



# Our goal

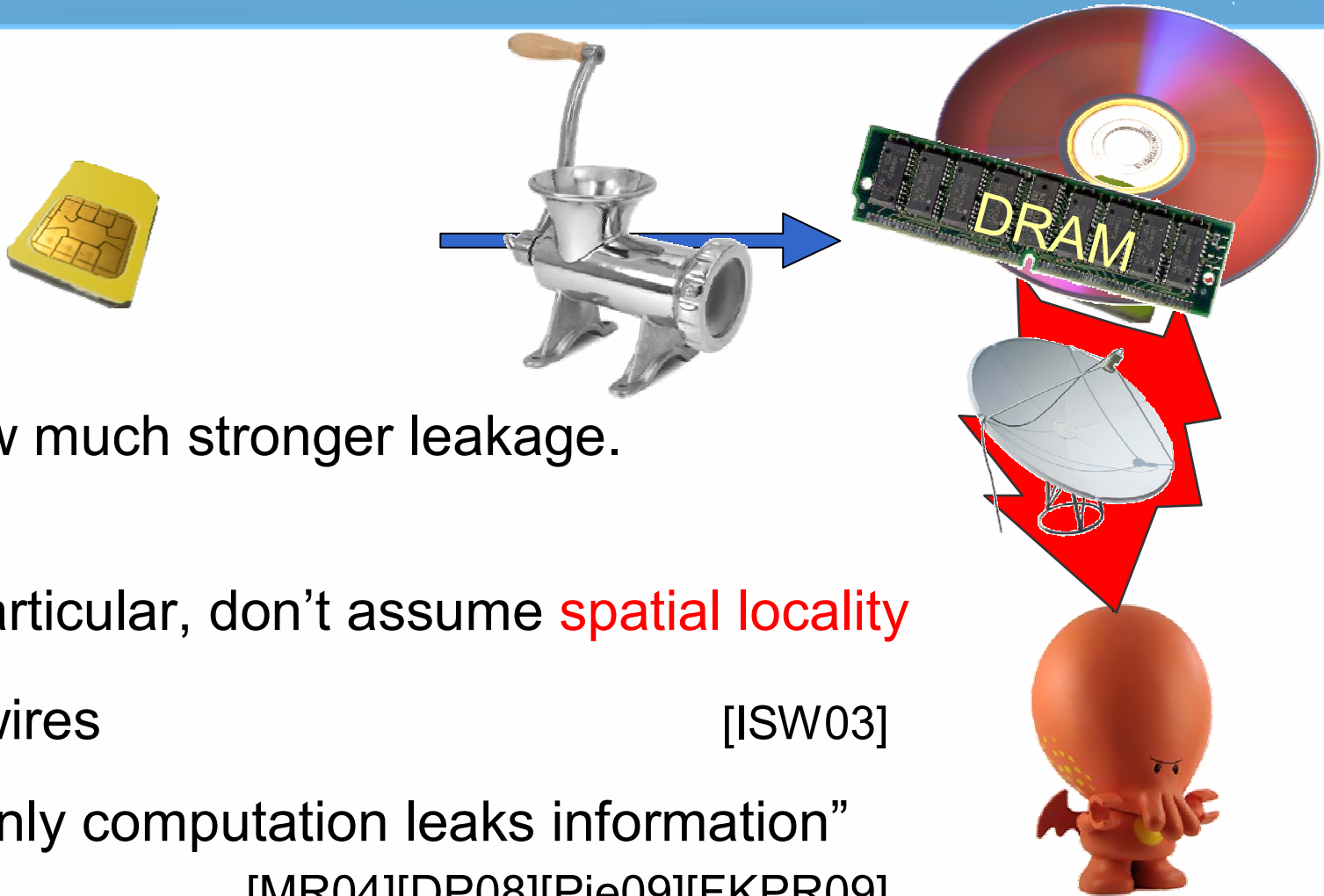


Allow much stronger leakage.

In particular, don't assume **spatial locality**

- $t$  wires [ISW03]
- “Only computation leaks information”  
[MR04][DP08][Pie09][FKPR09]

# Our goal



# Our main construction

A transformation that makes **any circuit** resilient against

- **Global adaptive leakage**

May depend on whole state and intermediate results, and chosen adaptively by a powerful on-line adversary.

- **Arbitrary total leakage**

Bounded just per observation.

[DP08]



But we must assume something:

- Leakage function is **computationally weak** [∈MR04]

- A simple **leak-free component** [∈MR04]

# Computationally-weak leakage

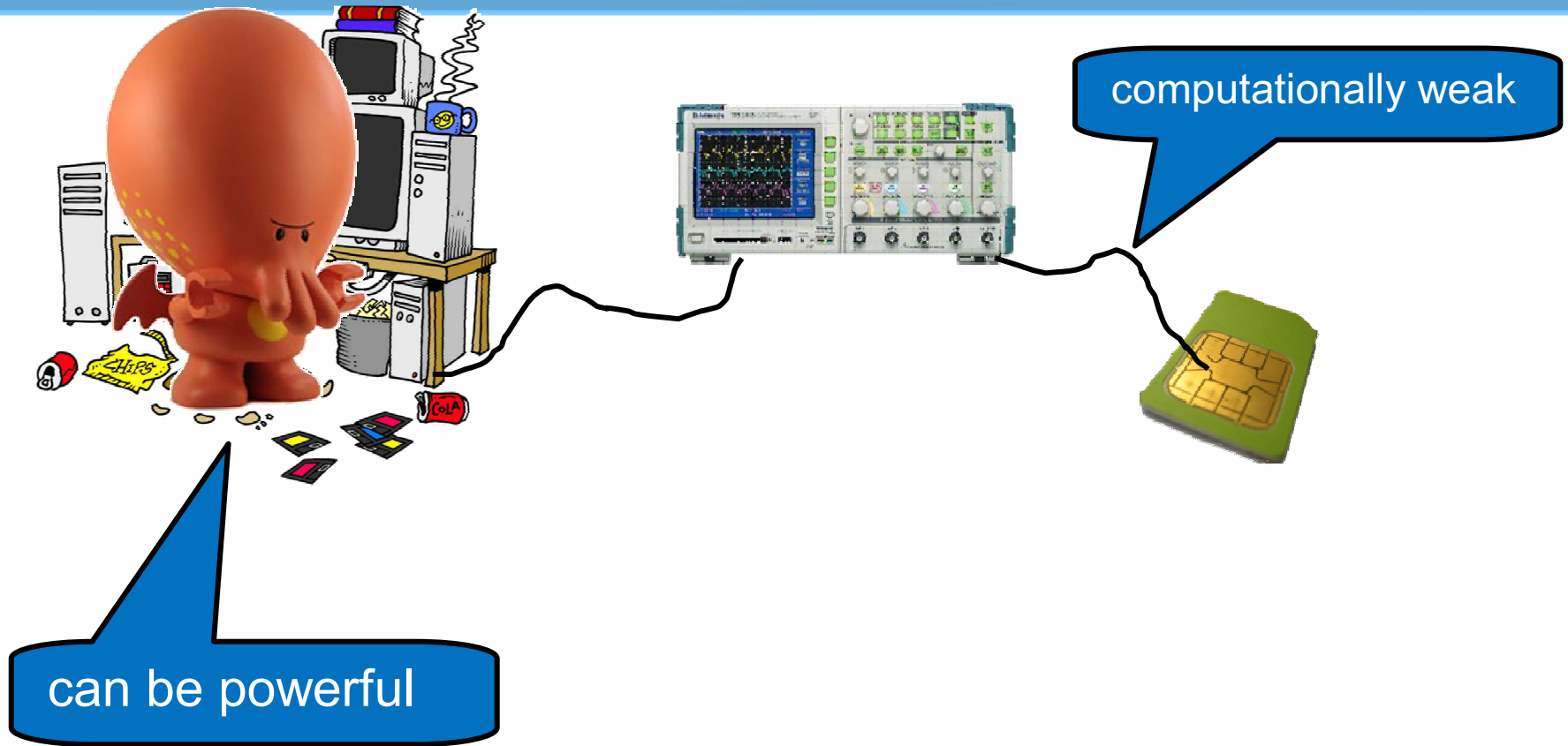


# Computationally-weak leakage

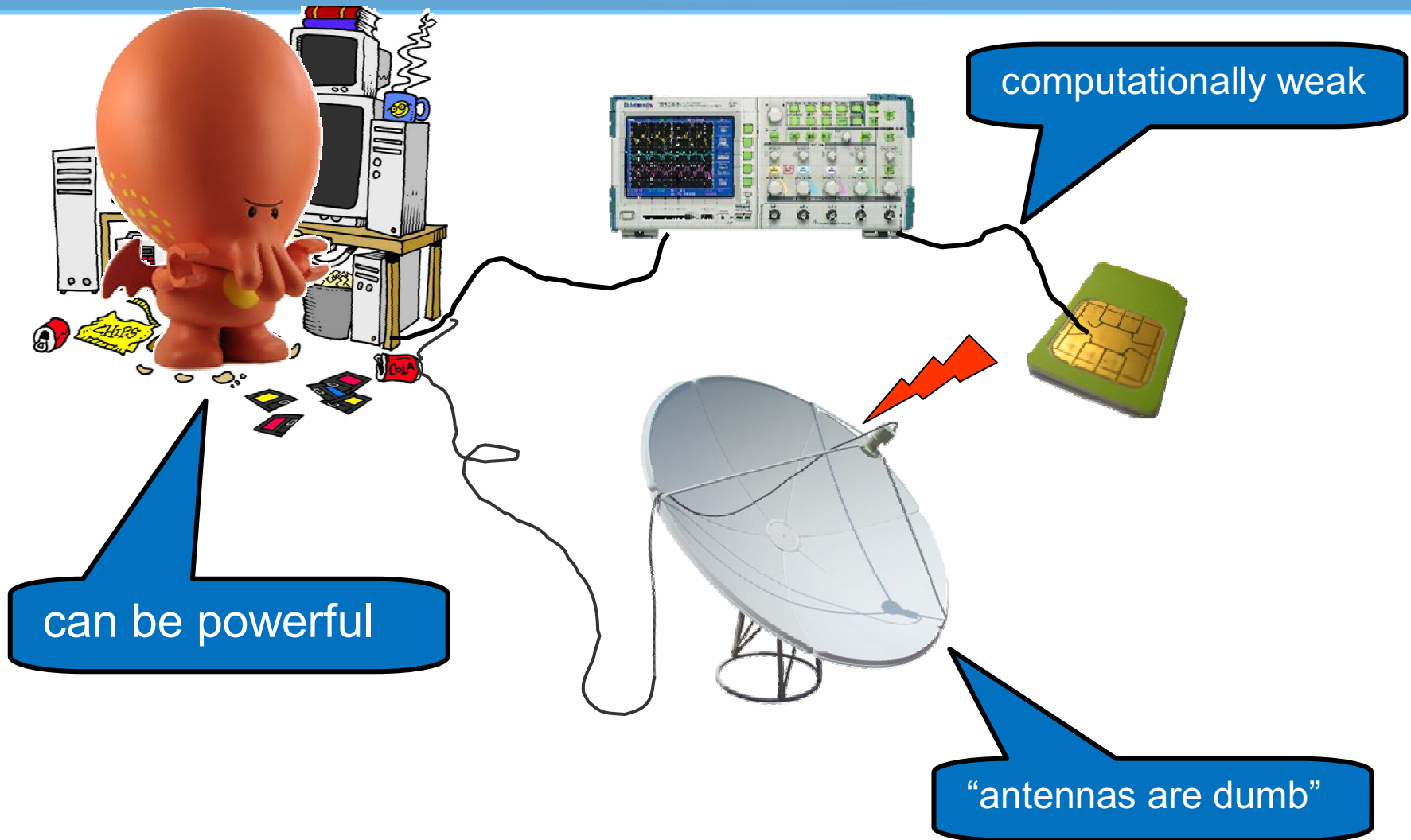


can be powerful

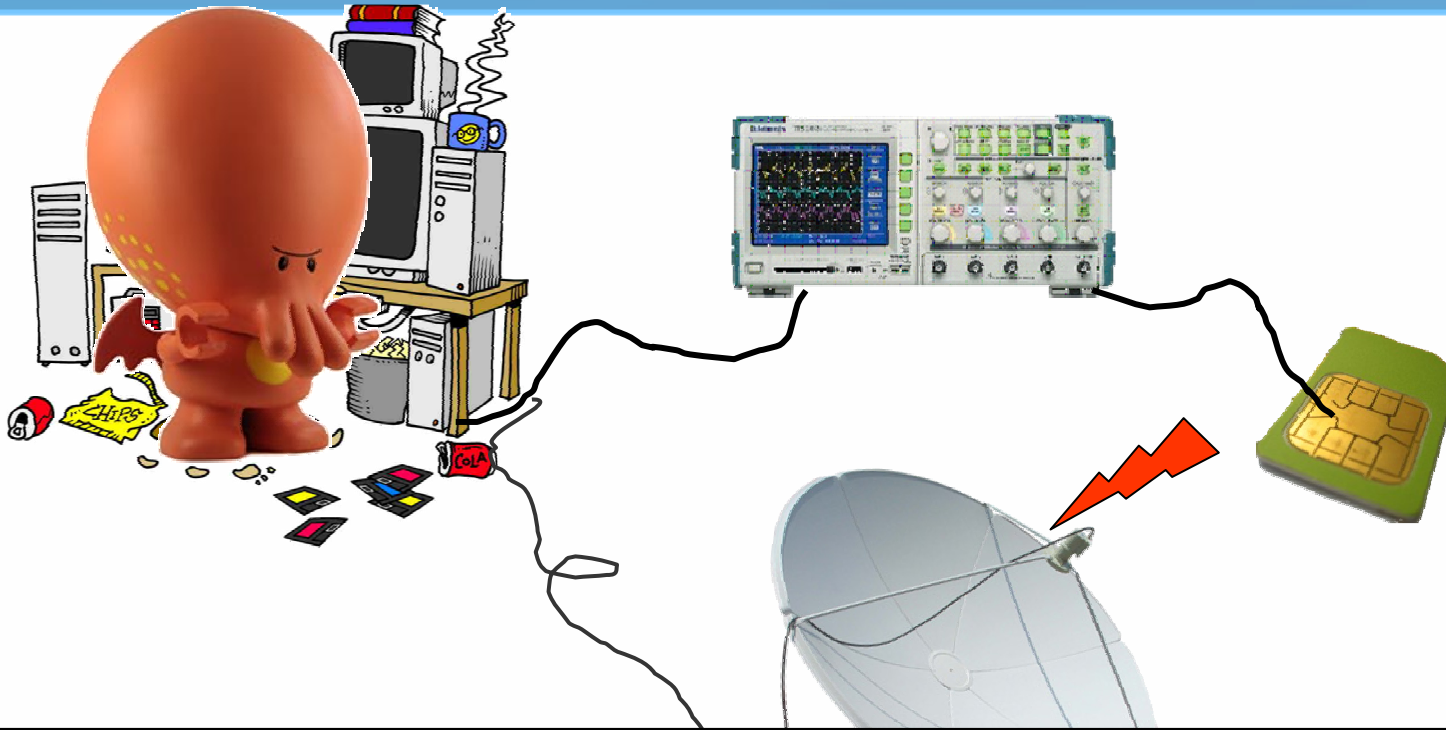
# Computationally-weak leakage



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# Computationally-weak leakage



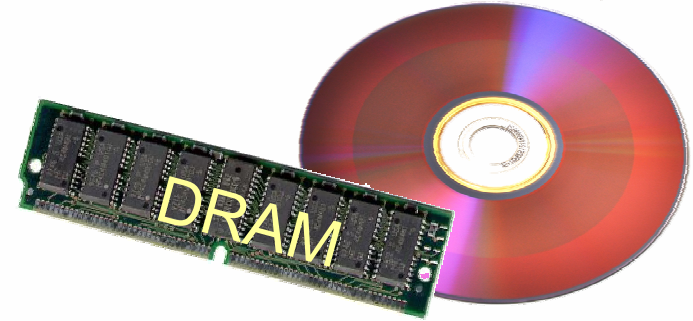
**Assumption: the observed leakage is a computationally-weak function of the device's internal wires.**



# Leak-free components

- **Secure memory**

[MR04][DP08][Pie09][FKPR09]

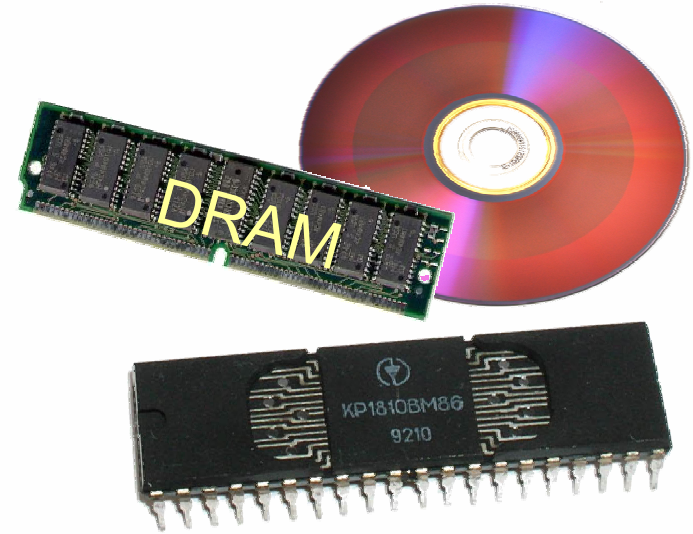


# Leak-free components

- **Secure memory**

[MR04][DP08][Pie09][FKPR09]

- **Secure processor** [G89][GO95]



# Leak-free components

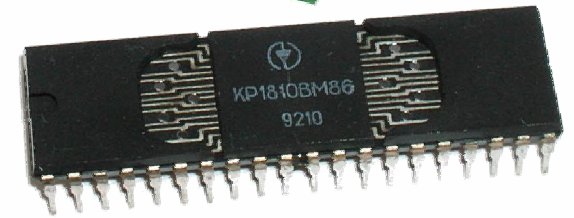
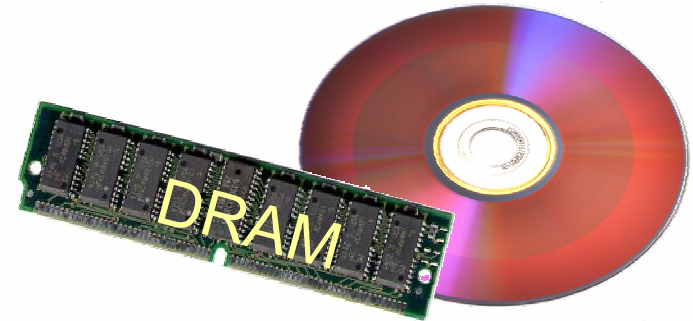
- **Secure memory**

[MR04][DP08][Pie09][FKPR09]

- **Secure processor** [G89][GO95]

- Here: simple component that samples from a fixed distribution, e.g: **draw strings with parity 0**.

- No stored secrets or state
- No input
  - consumable leak-free “tape roll”



# Results

- Constructions for **generic** circuit transformation using linear secret sharing schemes.
  - Example: unconditional security against  $AC^0$  leakage.
- Argue **necessity** of leak-free components (for “natural” constructions) by complexity-theoretic bounds/conjectures.
- General **proof technique** + additional applications.



<http://eprint.iacr.org/2009/341>